

INTRODUCTION TO ALGORITHMS

GIGA Doctoral School Introduction to Scientific Computing

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Outline

- Introduction
- Types of algorithms
- Classification of algorithms
- Expressing algorithms
- Constructs of an algorithm
- ► The concept of subalgorithm
- Examples
- Algorithm complexity



Introduction

Definition:

- An algorithm is step-by-step procedure with the aim of solving a problem.
- Algorithms are often used in many real life problems
- In computer science, an algorithm has a special meaning. It is defined to have these features:
 - An algorithm must have some data to operate on it
 - It must produce at least one result
 - It must terminate after a finite numbers of steps



Introduction

History:

- History of algorithms can be traced back to the ancient Greeks
- An efficient method for finding the Greatest Common Divisor was proposed by Euclid
- Study of algorithm was done by Mohammed ibn mussa al-Khowarizmi



Types of Algorithms

The types of algorithms depends on the type of **task** to be solved.

- * Searching
 - Designed to search for a given item in large data set
- * Sorting
 - Used to arrange data items in ascending or descending order
- * Compression
 - Meant to reduce the size of data and program files
 - Commonly used for compression of images, audio and video data



Types of Algorithms

* Fast Fourier Transforms

• Used in Digital Signal Processing (DSP)

* Encoding

• Used for encryption of data

* Geometric

- Used for identification of geometric shapes
- * Pattern Matching
 - Comparing images and shapes



Classification of Algorithms

Depending on the **strategy** used for **solving** a particular problem, algorithms are

classified as follows:

- Divide-and-Conquer Algorithms
 - A given problem is fragmented into sub-problems which are solved partially
 - The algorithm is stopped when further sub-division cannot be performed
 - These algorithms are frequently used in searching and sorting problems



Classification of Algorithms

Iterative Algorithms

- Certain steps are repeated in loops, until the goal is achieved
- An example of an iterative algorithm is sorting an array
- Greedy Algorithms
 - In a Greedy algorithm an immediately available best solution at each step is chosen
 - Useful for solving graph theory





Classification of Algorithms

- Back-Tracking Algorithms
 - In back tracking algorithms, all possible solutions are explored until the end is reached, afterwards the steps are traced back
 - These are useful in graph theory.
 - Back tracking algorithms are used frequently for traversing trees



- Describing algorithms requires a **notation** for expressing a **sequence** of steps to be performed.
- Algorithms can be expressed in many kinds of notation, including <u>natural languages</u>, <u>pseudocode</u>, <u>flowcharts</u>

Natural Language

- English words and sentences can be used to express statements and processing steps
 - For example, words like read, write, compute and set can be used for Input-Output operations, computations and assigning values to variables.
 - Comparison operations are expressed as equal to, less than, greater than
 - Arithmetical operations are expressed using words like add, subtract, divide and multiply
 - Control structures are expressed using sentences like repeat for, while, if, halt, exit
- Example: Find the largest element in a list/array of five integers.



What you would do?





What does it mean in natural language?

Input List

15

FindTheLargest





Could you express it in a more simple way?

FindTheLargest





Input/read: list of N integers

Set Largest to 0

Repeat the following N times

If the current number is greater than Largest, Set Largest to the current number

Output Largest

End



Use of Pseudocode

- Algorithms in natural language tend to be wordy and verbose
- Pseudocode provides an alternative way of expressing algorithms
- It is a mixture of natural language and programming notation
- In practice several conventions are used to write pseudocode

Input/read: list of N integers Set Largest to 0 Repeat the following N times If the current number is greater than Largest, Set Largest to the current number Output Largest End



Use of Pseudocode

- Algorithm is identified by a name
- Comments are enclosed in square brackets
- Assignment statement is coded using left arrow
- Operators : (+, -, *, /, <, >, =, !=)
- Input and Output : read and write
- Control Structures : if-then, if-then-else
- Repetitive operations : Repeat, for, while, until

FindTheLargest

- Input: A list of positive integers
- 1. Set Largest to 0
- 2. while (more integers)
- 3. if (the current integer is greater than Largest)
- 4. then
- 5. Set Largest to the value of the current integer
- 6. end if
- 7. End while
- 8. Return Largest
- 9. End

Flowchart



Flowchart Rules:

- 1. Flowchart is generally drawn from top to bottom
- 2. All boxes of flowchart must be connected
- 3. All flowchart start with terminal or process symbol
- Decision symbol have 2 exit points, one for YES (TRUE) and another for NO (FALSE)

Constructs of an algorithm

FindTheLargest

Input: A list of positive integers

- 1. Set Largest to 0
- 2. while (more integers)
- 3. if (the current integer is greater than Largest)
- 4. then
- 5.

Set Largest to the value of the current integer

- 6. end if
- 7. End while
- 8. Return Largest

9. End

do action 1 do action 2 . . . Sequence do action n if a condition is true. Then Decision do a series of actions Else do a series of actions While a condition is true. do action 2 Repetition do action n



Constructs of an algorithm

Constructs & pseudocode





Constructs of an algorithm

Constructs & Flowcharts





The concept of subalgorithm

FindTheLargest



3. Return Largest

End



Summation/Multiplication



Summation

Input: A list of integers

- 1. Set Sum to 0
- 2. While(more integers)
 - 2.1. Add current number to sum

End of while

3. Return Sum

End

Multiplication

Input: A list of integers

- 1. Set product to 1
- 2. While(more integers)
 - 2.1. Multiply current number by product

End of while

- 3. Return product
 - End



Sorting algorithms

Given a list, put it into some order

Input: sequence $(a_1, a_2, ..., a_n)$ of numbers. Output: permutation $(a'_1, a'_2, ..., a'_n)$ such that $a'_1 \le a'_2, \le ... \le a'_n$.



- □ We will see three types
 - Insertion sort
 - Selection sort
 - Bubble sort



Sorting algorithms

Insertion-Sort

It starts with a list with one element, and inserts new elements into their proper place in the sorted part of the list.





Sorting algorithms

Insertion-Sort





Flowchart

Examples of algorithms

Sorting algorithms

Insertion-Sort Pseudocode

Input: A list of integers (a₁, a₂,..., a_n)

- 1. for j = 2 to A.length
- 2. value = A[j]
- 3. Insert A[j] into the sorted sequence A[1..j-1]
- 4. i = j-1
- 5. While(i > 0 and A[i] > value)
- 6. A[i+1] = A[i]
- 7. i = i -1
- 8. End of while
- 9. A[i+1] = value
- 10. End of for
 - End (a'₁, a'₂, ..., a'_n) are sorted





Sorting algorithms

Selection-Sort

Find the smallest element in the unsorted list and swap it with the first element of the unsorted list.





Sorting algorithms





Sorting algorithms

Selection-Sort

	<u>Selection-Sort</u>
	Input: A list of integers (a ₁ , a ₂ ,, a _n)
1.	for i = 1 to A.length -1
2.	min = i
3.	/* check the element to be minimum */
4.	for j = i+1 to A.length
5.	if A[j] > A[min] then
6.	Min = j
7.	end if
8.	end for
9.	/* swap the minimum element with the current element */
10	. If indexMin != i then
11.	swap A[min] and A[i]
12.	end if
13.	End for
	End $(a'_{1}, a'_{2},, a'_{n})$ are sorted





Sorting algorithms

Bubble-Sort

- One of the least efficient algorithms
- It takes successive elements and « bubbles » them up/down in the list.





Sorting algorithms







Sorting algorithms

Bubble-Sort

Input: A list of integers	(a ₁ ,	a ₂ ,,	a _n
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- 1. for i = 1 to A.length
- 2. swapped = false
- 3. for j = 1 to A.length
- 4. [compare the adjacent elements]
- 5. **if** A[j] > A[j+1] then
- 6. [swap them]
- 7. swap(A[j], A[j+1])
- 8. swapped = true

→9.
10.
11.
12.
13.
14.
15.
16

end if

- end for
- [if no number was swapped that means
- list is sorted now, break the loop.]
- .3. if(not swapped) then
- 4. break
- 15. end if
- 16 End for
- 17. End $(a'_1, a'_2, ..., a'_n)$ are sorted



Searching algorithms

- Given a list, find a specific element in the list
- We will see two types
 - Linear search (sequential search)
 - Binary search



Searching algorithms





Linear search

Linear search running time

- How long does this take?
- If the list has n elements , worst case scenario is that it takes n « steps »
- Here, a step is considered a single step through the list



Searching algorithms



Binary search = List MUST be sorted!

Binary search running time

- How long does this take (worst case)?
- If the list has 8 elements
 - It takes 3 steps
- If the list has 16 elements
 - It takes 4 steps
- If the list has 64 elements
 - It takes 6 steps
- If the list has n elements
 - It takes log₂(n) steps





Space complexity

□ How much space is required?

Time complexity

□ How much time does it take to run algorithm?

Often, we deal with estimates!



Space complexity

Space complexity S(p) of an algorithm is the **total space** in memory taken by the algorithm to complete its execution with respect to the input size

S(p) = CONSTANT SPACE + AUXILARY SPACE

Constant space : is the space fixed for that algorithm, generally equals to space used by input and local variables

Auxilary space : is the extra/temporary space used by an algorithm

ONLY THE AUXILARY PART SHOULD BE CONSIDERED

S(p) = C + S(auxilary) = S(auxilary)

Space complexity

Summation

<u>Summation</u>	Input (a, n)
Input: a, b, c	Sum = 0
return a + b + c	for i in range (n)
End	sum = sum + a[i]
$S(p) = 1 + 1 + 1 = 3 \rightarrow No Auxiliary$	end for
	return Sum
	End

 $S(p) = (n*1+1+1) + 1 = n + 1 \rightarrow Auxilay = 1$



Time complexity



We analyze time complexity only for :

a) Very large input-size

b) Worst case scenario

Time complexity of an algorithm signifies the total **time** required by the program to run till its completion.

The time complexity of algorithms is most commonly expressed using the **Big O notation**.

Big O notation gives an **uper bound** of the complexity in the **worst** case, helping to quantify performance as the input size becomes **arbitrarily large**.

Time complexity



Big O notation

n: the size of the input Complexities ordered from smallest to largest Constant Time: O(1) Linear Time: O(n) Quadratic Time: O(n²) Cubic Time: O(n³)



Time complexity

Big O properties:

T(n) is a function describing the running time of a particular algorithm for an input of size n: T(n) = $n^3 + 3n^2 + 4n + 7$ T(n) $\approx n^3$ (n $\rightarrow \infty$)

≈ c n³ = *O(n³)*

Rule 1: a) Lower order terms should not be considered b) Constant multiplier should not be considered Example: $T(n) = 17 n^4 + 3 + 4n + 8 = O(n^4)$



Time complexity

Big O properties:	<u>Rule</u>: Running Time = ∑ Running Time o	of all fragments
int a;	For i = 0 to n; //simple statements	for (i = 0 ; i <n ;="" i++)<="" td=""></n>
a = 5		{ for (j = 0; j <n; j++)<="" td=""></n;>
a++;	Simple loop Fragment 2	{ //simple statements }
Simple statements Fragment 1 O(1)	O(n)) nested loop Fragment 3 O(n ²)



Time complexity



T(n) = O(1) + O(n)

or

 $T(n) = O(1) + O(n) + O(n^2) \approx O(n^2)$

<u>Rule</u>: Conditional Statements: Pick complexity of condition which is worst case





Take-home messages

- Algorithm is a step-by-step procedure to solve problems
- The types of algorithms depends on the type of **task** to be solved.
- Algorithms are classified based on the **strategy** used for **solving** problems.
- Algorithms can be expressed in : natural languages, pseudocode, and flowcharts.
- In one algorithm you could call another algorithm "concept of **subalgorithm**".
- Algorithm complexity is seen as Space complexity and time complexity.







Thank you for your attention

